## 33-658 Quantum Computing and Quantum Information Homework 8

- 1. (i) prove  $\mathbf{C}_{ij} = \tilde{\mathbf{n}}_i + \mathbf{n}_i \mathbf{X}_j$ . (ii) Apply algebraic manipulations (i.e. do not use matrix arithmetic) to prove:  $\mathbf{S}_{ij} = (1/2)(\mathbf{1} + \mathbf{X}_i \mathbf{X}_j + \mathbf{Y}_i \mathbf{Y}_j + \mathbf{Z}_i \mathbf{Z}_j) = (1/2)(1 + \vec{\sigma}_i \cdot \vec{\sigma}_j)$ .
- 2. Prove the identities





$$a \cdot x = a_0 x_0 \oplus a_1 x_1 \oplus \dots \oplus a_{n-1} x_{n-1}$$

is the modulo 2 sum of products of bits. Suppose that we are given a quantum circuit that evaluates f(x) for any input x but that we do not know the value of a. To determine a on a classical computer would require n evaluations of f. This exercise, which is based on the analysis in Mermin, shows how to determine a with a *single* evaluation of f on a quantum computer. (i) Let  $\mathbf{U}_f$  be an f-controlled unitary transformation on the tensor product of the n-bit input register initially in state  $|x\rangle_n$  and a 1-bit output register initially in state  $|y\rangle_1 = \mathbf{HX}|0\rangle = (|0\rangle - |1\rangle)/\sqrt{2}$ . Show that

$$\mathbf{U}_f |x\rangle_n |y\rangle_1 = (-1)^{f(x)} |x\rangle_n |y\rangle_1$$

(ii) Show that

$$\mathbf{H}^{\otimes n} |x\rangle_n = \frac{1}{2^{n/2}} \sum_{y=0}^{2^n - 1} (-1)^{x \cdot y} |y\rangle_n$$

(iii) Show that

$$(\mathbf{H}^{\otimes n} \otimes \mathbf{H}) \mathbf{U}_f (\mathbf{H}^{\otimes n} \otimes \mathbf{H}) |0\rangle_n |1\rangle_1 = |a\rangle_n |1\rangle_1.$$

Note that the unknown value of a appears in the *input* register!